## Theory and Explicit Design of a Path Planner for an SE(3) Robot

Zhaoqi Zhang, Yi-Jen Chiang, Chee Yap

Oct 7,2024 WAFR 2024, Chicago



### **Problem and Notations**

- We deal with kinematic path planning problem.
- Our robot is an isosceles right triangle  $\mathcal{AOB}$  in  $\mathbb{R}^3$  (Delta robot).
  - We call the area in physical space possessed by the robot under a given configuration  $\gamma$  the **footprint** of  $\gamma$ , denoted by  $Fp(\gamma)$ .

### Input

- In physical space ℝ<sup>3</sup>, an obstacle set, denoted by Ω ⊆ ℝ<sup>3</sup>.
- Start and goal configurations  $\alpha$  and  $\beta$  in configuration space *Cspace*.
- Resolution parameter  $\varepsilon > 0$

### Output

- A **path** (continuous map) from  $\alpha$  to  $\beta$ .
- Or **NO-PATH**.







### **Resolution Exactness**

- We use SSS (Soft Subdivision Search) framework. The output of SSS framework is resolution exact, i.e.,
- There exists some K > 1 (independent of input), such that:
  - (P) if there is a path of clearance  $K\varepsilon$ , it returns a path;
  - (N) if there is no path of essential clearance  $\varepsilon/K$ , it returns NO-PATH.
- The SSS framework is currently the only complete method for path planning (other than exact computations) that does not have the halting problem.
- Resolution exactness of SSS is guaranteed by our Fundamental Theorem which depends on 5 axioms (see next).



4

## SSS Axioms (constants $\sigma$ , $D_0$ , $L_0$ , $C_0$ )

- (A0) Softness.
  - The predicate  $\tilde{C}$  is a soft classifier for *Cspace*;
  - The SSS is **effective** if the predicate is  $\sigma$ -effective.
- (A1) Bounded Expansion.
  - There is a **subdivision constant**  $D_0 \ge 1$  such that each box can be subdivided into at most  $D_0$  children and the aspect ratio of each box is no more than  $D_0$ .
- (A2) Lipschitz Clearance.
  - The footprint satisfies a Lipschitz constant  $L_0 > 0$ :

 $d_H(Fp(\gamma), Fp(\gamma')) \le L_0 d(\gamma, \gamma')$ 

where  $d_H$  is the Hausdorff distance in  $\mathbb{R}^k$ .

- (A3) Good Atlas.
  - The subdivision atlas has an atlas constant  $C_0 \ge 1$ .
- (A4) Translational Cell.
- **NYU** The boxes are **translational**.

## **Compare to Sampling Approach**

- Finding a path:
  - PRM/RRT/EST/SRT/etc.
  - Sampling functions, local planners, tree/graph-based planners [2].
  - The  $C_{free}$  must satisfy  $\varepsilon$ -goodness [3] and  $\delta$ -clearance [4].
- Checking NO-PATH (infeasibility proof):
  - Learn and validate Cobs manifold [5].
  - The  $C_{obs}$  must be entirely  $\varepsilon$ -blocked [5].
- Requires "promise input"



From Li and Dantam [5]:

### **Zero Problem**

7

- Consider a planar disc robot with radius 1 in  $\mathbb{R}^2$ .
  - The configuration space is  $\mathbb{R}^2$ .
  - Let the obstacle set be  $\Omega = \{(x, y): x < -1 \text{ or } x > 1, y = 0\}.$
  - The start configuration  $\alpha = (0, 1)$ , goal configuration  $\beta = (0, -1)$ .
- Neither  $C_{free}$  is  $\varepsilon$ -good, nor  $C_{obj}$  is  $\varepsilon$ -blocked.
- To determine if configuration (0,0) is free,
  - We must solve this zero problem.
    - $\{(x, y): x^2 + y^2 \le 1\} \cap \Omega = \emptyset$ ?
  - We must use exact computation.
    - In general, it is at least single exponential time in the degree of freedom.
- SSS planner will return NO-PATH.

 $\begin{array}{c} c_{spuce} \\ c_{free} \\ \hline \\ c_{obj} \\ \hline \\ c_{obj} \\ \hline \end{array}$ 

Cspace

### **Predicates in SSS framework**

 A predicate C classifies each configuration box B into FREE/MIXED/STUCK (a.k.a. empty/mixed/full) :

$$C(B) = \begin{cases} \mathsf{FREE} & \forall \gamma \in B, \gamma \in C_{free} \\ \mathsf{STUCK} & \forall \gamma \in B, \gamma \notin C_{free} \\ \mathsf{MIXED} & otherwise \end{cases}$$

- A soft predicate  $\tilde{C}$  is used for implementations that gives weaker but correct classifications.
  - Conservative:

 $\tilde{C}(B) \neq \text{MIXED} \text{ implies } C(B) = \tilde{C}(B);$ 

- Convergent:
  - If  $\{B_i\}$  is a sequence of boxes such that  $B_{i+1} \subseteq B_i$  and  $\bigcap_{i=1}^{\infty} B_i = \{p\}$  for some  $p \in Cspace$ , then

 $\tilde{C}(B_i) = C(p)$  for *i* large enough.



### **SSS Framework**

### • Priority queue Q:

- Controls the search of MIXED boxes.
- GetNext() may adopt different strategies.
- Find:
  - Union find method preserves connected components.

### • Expand:

- Subdivide boxes into subboxes;
- Classify each children:
  - If FREE, then add into the Union Find;
  - If MIXED, then add into the Q;



### SSS Framework

Input: Start configuration  $\alpha$ , goal configuration  $\beta$ , obstacle  $\Omega$ , resolution parameter  $\varepsilon$ . Output: A path  $\bar{P}$  or NO-PATH.

$$\begin{split} 1. &\rhd Initialization \\ & \text{While } (\widetilde{C}(\text{Box}(\alpha)) \neq \text{FREE}), \\ & \text{if } l(\text{Box}(\alpha)) < \varepsilon, \text{ return NO-PATH;} \\ & \text{else, } \texttt{Expand}(\text{Box}(\alpha)). \\ & \text{While } (\widetilde{C}(\text{Box}(\beta)) \neq \text{FREE}), \\ & \text{if } l(\text{Box}(\beta)) < \varepsilon, \text{ return NO-PATH;} \\ & \text{else, } \texttt{Expand}(\text{Box}(\beta)). \end{split}$$

### 2. $\triangleright$ Main Loop

While  $(\text{Find}(\text{Box}(\alpha)) \neq \text{Find}(\text{Box}(\beta)),$ if Q is empty, return NO-PATH  $B \leftarrow Q.\text{GetNext}()$ Expand(B).

3.  $\triangleright$  Search

Compute a FREE channel P from  $Box(\alpha)$  to  $Box(\beta)$ Generate and return the canonical path  $\bar{P}$  inside P.



FREE

MIXED

STUCK

*ɛ*-small

### **Subdivision Process**



- Box space:
  - This is a correspondence from  $\mathbb{R}^7$  to the configuration space SE(3).
  - The configuration space is  $SE(3) \cong \mathbb{R}^3 \times SO(3)$ .
    - Boxes in  $\mathbb{R}^3$  are **translational** boxes.
    - Boxes in SO(3) are **rotational** boxes (embedded into  $\mathbb{R}^4$ ).
- Subdivide and classify:
  - Expand the boxes containing  $\alpha$  and  $\beta$  until they are contained in FREE boxes;
  - Expand the "next" box in Q;
  - Stop when the boxes containing  $\alpha$  and  $\beta$  are in the same connected components, or the *Q* is empty.
- Build a FREE channel from the connected components of  $\alpha$  and  $\beta$ .

## **Approximate Footprint**

- The soft predicate will be given by an approximate footprint  $\widetilde{Fp}$ .
  - The obstacle set  $\Omega$  will be inputted as a set of **features**  $\Phi$  consists of points, edges, triangles and polyhedrons.
  - The soft predicate is defined as

$$\widetilde{C}(B) = \begin{cases} \mathsf{FREE} & \widetilde{Fp}(B) \land \Phi = \emptyset \text{ and } \widetilde{Fp}(B) \not\subseteq \Omega \\ \mathsf{STUCK} & \widetilde{Fp}(B) \land \Phi = \emptyset \text{ and } \widetilde{Fp}(B) \subseteq \Omega \\ \mathsf{MIXED} & \widetilde{Fp}(B) \land \Phi \neq \emptyset \end{cases}$$

• To make the soft predicate conservative and convergent, the approximate footprint satisfies : there is some  $\sigma > 1$  such that

 $\widetilde{Fp}(B/\sigma) \subseteq Fp(B) \subseteq \widetilde{Fp}(B)$ 



• This property is called  $\sigma$ -effectivity.

### **Approximate Footprint for Delta Robot**



NYU

### **Apply to Delta robot**

• Regard the approximate footprint as the union of 4 fat sets:



• Turn into Parametric Separation Queries:



**NYU** Sep(\*, f) > r(B)?

Sep(\*, f) > r(B)?

Sep(\*, f) > d(B)?

0.4 0.6 0.8

Sep(\*, f) > 0?

### **Performance (Very preliminary)**



RGB - xyz

This environment find a path with 2996 boxes in 9.52297s.



RGB - xyz

This environment find a path with 8642 boxes in 20.9203s.



RGB - xyz

This environment find NO-PATH with 1866 boxes in 2.90841s.

### Conclusion

- This is the first explicit complete SE(3) path planner.
  - Explicit:
    - No invocation of an optimizer.
    - No Newton iteration.
    - No machine learning.
  - All computations are reduced to semi-algebraic tests.
- A full-scale implementation will require additional search techniques (on-going work).



### Reference

[1] C. Wang, Y.-J. Chiang, and C. Yap. On soft predicates in subdivision motion planning. *Comput. Geometry: Theory and Appl. (Special Issue for SoCG'13),* 48(8):589–605, Sept. 2015.

[2] Andreas Orthey and Constantinos Chamzas and Lydia E. Kavraki. Sampling-Based Motion Planning: A Comparative Review. *Annual Reviews.* Vol. 7:285-310, Nov. 2023.

[3] Kavraki, L.E., Latombe, J.C., Motwani, R., Raghavan, P. Randomized query processing in robot path planning. JCSS 57(1), 50–60 (1998)

[4] Karaman, S., Frazzoli, E.: Sampling-based algorithms for optimal motion planning. IJRR 30(7), 846–894 (2011)

[5] Sihui Li and Neil T. Dantam. Exponential Convergence of Infeasibility Proofs for Kinematic Motion Planning. WAFR 22, 294–311 (2023)



# Thanks for Listening!



